

# THE DISTANCE FOR THE BÉZIER CURVES AND DEGREE REDUCTION

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ABSTRACT. An algorithmic approach to degree reduction of Bézier curves is presented. The algorithm is based on the matrix representations of the degree elevation and degree reduction processes. The control points of the approximation are obtained by the generalized least square method. The computations are carried out by minimizing the  $L_2$  and discrete  $l_2$  distance between the two curves.

## 1. INTRODUCTION

The Bézier curves are basically and widely used in CAGD – short for Computer Aided Geometric Design. The Bézier curves were independently developed by P. de Casteljaeu about 1959[2] and by P. Bézier about 1962[1]. The underlying mathematical theory is based on the concept of Bernstein polynomials. De Casteljaeu directly exploited this relationship; but it was not before 1970 that R. Forrest[11] discovered the connection between Bézier’s work and Bernstein polynomials. Bézier and de Casteljaeu developed their theories as part of CAD systems that were being built up at two French car companies, Renault and Citroën. The Renault system UNISURF (by Bézier) was soon described in several publications; this is the reason that the underlying theory now bears Bézier’s name. Bézier curves and surfaces are now established as the mathematical basis of many CAD systems, they have also become a major tool for the development of new methods for curve and surface descriptions. Farin[10] summarize the basic theory of such curves and provide many relevant references.

The Bézier representation uses Bernstein polynomials as basis functions for the linear space of polynomials. In terms of the Bernstein polynomials of degree  $n$ ,

$$B_k^n(t) = \binom{n}{k} (1-t)^{n-k} t^k, \quad 0 \leq t \leq 1, \quad k = 0, \dots, n,$$

a parametric polynomial curve of degree  $n$  ( $n > 0$ ) in the plane, can be expressed as

$$b^n(t) = \sum_{k=0}^n b_k B_k^n(t), \quad b_k \in \mathbb{R}^2.$$

The points  $b_k$ ,  $k = 0, \dots, n$  are called the *control points* for the polynomial, and the polygon formed by joining successive control points is the *control polygon*. Notice that  $b_0$  and  $b_n$  are the endpoints of the curve corresponding to  $t = 0$  and  $t = 1$ ; we shall refer to these particular points as *anchor points*. Moreover, the vector  $b_1 - b_0$  and  $b_n - b_{n-1}$  define the tangents to the curve at the two anchor points respectively.

In general degree reduction of Bézier curves address the following problem.

PROBLEM 1 (Degree Reduction). Let  $\{b_i\}_{i=0}^n \subset \mathbb{R}^2$  be a given set of control points which define the Bézier curve

$$b^n(t) = \sum_{i=0}^n b_i B_i^n(t), \quad 0 \leq t \leq 1$$

of degree  $n$ . Then find another point set  $\{c_i\}_{i=0}^m \subset \mathbb{R}^2$  defining the approximative Bézier curve

$$c^m(t) = \sum_{i=0}^m c_i B_i^m(t), \quad 0 \leq t \leq 1$$

of lower degree  $m < n$  so that a suitable distance function  $d(b^n, c^m)$  between  $b^n$  and  $c^m$  is minimized.

In the literatures([3], [4], [6], [7], [9], [11], [14], [15], [17], [19]) one can find several schemes producing solutions for this approximation problem. These schemes mainly differ in the choice of the distance function and requiring the solution to be either best or only nearly best relative to the distance function. For instance, one special type of degree reduction schemes works recursively by lowering the degree only by one in every step – a procedure commonly known as economization.

The examples for such a stepwise method was recently given in Eck[6] or Park, Choi and Kimn[19] where a very simple geometric construction of the new control points in each step is described. And the method allow the detailed error analysis for the other method (e.g. Forrest[11] and Farin[9], see Park and Choi[18]). However, this general construction contains some scalar – valued degrees of freedom which are then chosen in such a way that the maximal Euclidean distance

$$d_\infty(b^n, c^{n-1}) = \max_{0 \leq t \leq 1} \|b^n(t) - c^{n-1}(t)\|$$

between two curves with respect to the given parameterization is minimized.

The derivation is mainly based on the so-called constrained Chebyshev polynomials. Unfortunately, the constrained Chebyshev polynomials are not known explicitly so their coefficients have to be determined numerically, which itself needs a lot of implementation effort.

This major disadvantage is avoided in the current paper. In more detail, we minimize the least squares distance function

$$d_2(b^n, c^m) = \sqrt{\int_0^1 \|b^n(t) - c^m(t)\|^2 dt}.$$

The algorithm presented is faster, more stable and much easier to implement. Moreover, the procedure apply to reduce the degree from  $n$  to  $m$  as only one step.

## 2. DEGREE ELEVATION AND $L_2$ DISTANCE

Suppose we were designing with Bézier curves trying to use a Bézier curve of degree  $n$ . After modifying the polygon a few times, it may turn out that a degree  $n$  curve does not possess sufficient flexibility to model the desired shape. One way to proceed in such situation is to increase the flexibility of the polygon by adding another vertex to it. As a first step, one might want to add another vertex yet leave the shape of the curve unchanged – this corresponds to raising the degree of the Bézier curve by one. We can show that new vertices  $b_i^{(1)}$  are obtained from the old polygon by piecewise linear interpolation at the parameter values  $i/(n+1)$

$$b_i^{(1)} = \frac{i}{n+1}b_{i-1} + \left(1 - \frac{i}{n+1}\right)b_i, i = 0, 1, \dots, n+1. \quad (1)$$

We can rewrite the formular (1) as a linear system  $T_n B = B^{(1)}$ , where the  $(n+2) \times (n+1)$  matrix  $T_n$  is

$$T_n = \frac{1}{n+1} \begin{pmatrix} n+1 & 0 & 0 & \dots & 0 & 0 & 0 \\ 1 & n & 0 & \dots & 0 & 0 & 0 \\ 0 & 2 & n-1 & \dots & 0 & 0 & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots & \vdots \\ 0 & 0 & 0 & \dots & n-1 & 2 & 0 \\ 0 & 0 & 0 & \dots & 0 & n & 1 \\ 0 & 0 & 0 & \dots & 0 & 0 & n+1 \end{pmatrix}$$

and the  $(n+1)$  vector  $B$  and the  $(n+2)$  vector  $B^{(1)}$  are

$$\begin{aligned} B &= (b_0, b_1, \dots, b_n)^t \\ B^{(1)} &= (b_0^{(1)}, b_1^{(1)}, \dots, b_n^{(1)})^t. \end{aligned}$$

We may repeat this process and then obtain a sequence of controls points. After  $r$  degree elevations, we have a linear system  $T_{n,r} B = B^{(r)}$ , where the  $(n+r+1) \times (n+1)$  matrix

$$T_{n,r} = T_{n+r-1} T_{n+r-2} \dots T_{n+1} T_n$$

has elements

$$t_{i+j,i} = \frac{\binom{n}{i} \binom{r}{j}}{\binom{n+r}{i+j}}, \quad \begin{cases} i = 0, 1, \dots, n \\ j = 0, 1, \dots, r \end{cases}$$

The sum of any row and any column of the matrix  $T_{n,r}$  are 1 and  $\frac{n+r+1}{n+1}$  respectively, i.e. for any  $i$ ,

$$\sum_{k=0}^{n+r} t_{i,k} = 1,$$

and for any  $k$ ,

$$\sum_{i=0}^n t_{i,k} = \frac{n+r+1}{n+1}.$$

For the degree reduction of any given curves, we must compute a distance of two Bézier curves. The most appropriate metric in geometrical terms would be the *Hausdorff distance*[5]. Suppose  $(M, d)$  is a metric space with subsets  $A$  and  $B$ . We define the Hausdorff metric  $d_H$  by

$$d_H(A, B) = \max \left\{ \sup_{x \in A} d(x, B), \sup_{y \in B} d(y, A) \right\},$$

where

$$d(x, B) = \inf_{y \in B} d(x, y).$$

If we regard a planar curve as simply a locus of points without any underlying parameterization, the Hausdorff metric for two such curves is essentially the radius of the largest circle with its center on one curve and touching the other curve. For general parametric curves, this measure is truly independent of the relative parameterizations of two curves. Emery[8] presents a method for explicit computation of Hausdorff metric for piecewise

linear curves, but the computation of Hausdorff distance  $d_H$  of two nonlinear curves is not so easy. So we define and use the  $L_2$  distance for the Bézier curves.

We first consider the functional case of the Bézier curves for computation of  $L_2$  distance of the two Bézier curves. Let  $a^n$  and  $b^m$  be the functional Bézier curves of degree  $n$  and  $m$  ( $m < n$ ), i.e.

$$a^n(t) = \sum_{k=0}^n a_k B_k^n(t), \quad b^m(t) = \sum_{k=0}^m b_k B_k^m(t),$$

where the coefficients  $a_k$  and  $b_k$  are real numbers. The  $L_2$  distance of the two Bézier curves  $a^n$  and  $b^m$  is defined as following:

$$\begin{aligned} d_2(a^n, b^m) &= \left\{ \int_0^1 |a^n(t) - b^m(t)|^2 dt \right\}^{\frac{1}{2}} \\ &= \left\{ \int_0^1 \left| \sum_{k=0}^n a_k B_k^n(t) - \sum_{k=0}^m b_k B_k^m(t) \right|^2 dt \right\}^{\frac{1}{2}} \end{aligned}$$

Using the matrix  $T_{m,r}$ , we can elevate the degree of  $b^m$  from  $m$  to  $n$ ,

$$B^{(r)} = T_{m,r} B.$$

Then, the curve  $b^m$  of degree  $m$  is re-written as a curve of degree  $n$

$$b^m(t) = b^{(r)}(t) = \sum_{k=0}^n b_k^{(r)} B_k^n(t),$$

and the distance is

$$\begin{aligned} d_2(a^n, b^m) &= d_2(a^n, b^{(r)}) \\ &= \left\{ \int_0^1 \left| \sum_{k=0}^n a_k B_k^n(t) - \sum_{k=0}^n b_k^{(r)} B_k^n(t) \right|^2 dt \right\}^{\frac{1}{2}} \\ &= \left\{ \int_0^1 \left| \sum_{k=0}^n (a_k - b_k^{(r)}) B_k^n(t) \right|^2 dt \right\}^{\frac{1}{2}}. \end{aligned}$$

Let  $c_k = a_k - b_k^{(r)}$  for all  $k$ . So, we compute the  $L_2$  norm of a (functional) Bézier curve of degree  $n$ .

The product of Bernstein polynomials is

$$B_i^n(t) B_j^m(t) = \frac{\binom{n}{i} \binom{m}{j}}{\binom{n+m}{i+j}} B_{i+j}^{n+m}(t) \quad (2)$$

and the integration is

$$\int_0^1 B_k^n(t) dt = \frac{1}{n+1}. \quad (3)$$

From these equations (2) and (3), we obtain the following computation for the  $L_2$  norm of the functional Bézier curve  $c^n$ :

$$\begin{aligned} \|c^n\|_2^2 &= \int_0^1 \left| \sum_{k=0}^n c_k B_k^n(t) \right|^2 dt \\ &= \int_0^1 \sum_{i,j} c_i c_j B_i^n(t) B_j^n(t) dt \\ &= \sum_{i,j} c_i c_j \int_0^1 \frac{\binom{n}{i} \binom{n}{j}}{\binom{2n}{i+j}} B_{i+j}^{2n}(t) dt \\ &= \frac{1}{2n+1} \sum_{i,j} c_i c_j \frac{\binom{n}{i} \binom{n}{j}}{\binom{2n}{i+j}} \end{aligned}$$

Let  $(n+1) \times (n+1)$  matrix  $Q_n$  be the

$$Q_n = \frac{1}{2n+1} \left[ \frac{\binom{n}{i} \binom{n}{j}}{\binom{2n}{i+j}} \right]_{(n+1) \times (n+1)}$$

Then,  $L_2$  norm of the Bézier curve  $c^n$  is

$$\|c^n\|_2^2 = C^t Q_n C.$$

The matrix  $Q_n$  is a real symmetric matrix. The following lemma tell us the equivalent conditions for the real symmetric matrix to be positive definite matrix[16].

**LEMMA 1.** *Each of the following tests is a necessary and sufficient condition for the real symmetric matrix  $A$  to be positive definite:*

1.  $x^t A x > 0$  for all nonzero vectors  $x$ .
2. All the eigenvalues of  $A$  are positive.
3. All the upper left submatrices have positive determinants.

From the definition of the matrix  $Q_n$  and the mathematical induction, all the upper left submatrices of the matrix  $Q_n$  have positive determinants. Applying the Lemma 1, the matrix  $Q_n$  is a real symmetric positive definite.

Thus, we obtain the following theorem for the  $L_2$  distance between the Bézier curves  $a^n$  of degree  $n$  and the Bézier curve  $b^m$  of degree  $m$ .

**THEOREM 2.** *The  $L_2$  distance between the two Bézier curves  $a^n$  and  $b^m$  is*

$$d_2(a^n, b^m) = d_2(a^n, b^{(r)}) = \sqrt{D^t Q_n D}, \quad (4)$$

where  $D = A - T_{m,r} B$  and  $A = (a_0, \dots, a_n)^t$  and  $B = (b_0, \dots, b_m)^t$ .

The sum of any row and any column of the matrix  $Q_n$  are both equal to  $\frac{1}{n+1}$  i.e.

$$\sum_{j=0}^n q_{i,j} = \sum_{i=0}^n q_{i,j} = \frac{1}{n+1}.$$

### 3. DEGREE REDUCTION

Replacing the distance function  $d$  in the Problem 1 and the distance  $d_2$  as (4), then the Problem 1 rewritten as the following:

PROBLEM 2 ( $L_2$  Degree Reduction). Find another points set  $\{c_i\}_{i=0}^m$  so that the least squares distance

$$d_2(b^n, c^m) = d_2(b^n, c^{(r)}) = \sqrt{D^t Q_n D}$$

between  $\{b_i\}_{i=0}^n$  and  $\{c_i^{(r)}\}_{i=0}^n$  is minimized.

Note that  $T_{m,r}C = C^{(r)}$  and  $B - C^{(r)} = D$  where  $D = (d_0, d_1, \dots, d_n)^t$ ,  $d_i = b_i - c_i^{(r)}$ ,  $i = 0, 1, \dots, n$ .

For developing the method rewrite the  $D^t Q_n D$ ,

$$\begin{aligned} D^t Q_n D &= [B - C^{(r)}]^t Q_n [B - C^{(r)}] \\ &= [B - T_{m,r}C]^t Q_n [B - T_{m,r}C] \\ &= B^t Q_n B - 2C^t T_{m,r}^t Q_n B + C^t T_{m,r}^t Q_n T_{m,r} C. \end{aligned}$$

One method of obtaining the vector  $C$  is so-called method of least squares. This method consists of minimizing  $D^t Q_n D$  with respect to  $C$ . Choosing the vector  $\hat{C}$  that value of  $C$  which minimize  $D^t Q_n D$  involves differentiating  $D^t Q_n D$  with respect to the elements of  $C$ . Equating  $\partial(D^t Q_n D)/\partial C$  to zero and writing the resulting equations in terms of  $\hat{C}$ , we find that these equations are

$$T_{m,r}^t Q_n T_{m,r} \hat{C} = T_{m,r}^t Q_n B.$$

They are known as the normal equations[12].

From the definition of the matrix  $T_{m,r}$  and  $Q_n$ , we have that the matrix product  $T_{m,r}^t Q_n T_{m,r}$  is  $Q_m$  i.e.

$$T_{m,r}^t Q_n T_{m,r} = Q_m.$$

Hence, the real symmetric positive definite matrix  $T_{m,r}^t Q_n T_{m,r} = Q_m$  is invertible. Provided  $(T_{m,r}^t Q_n T_{m,r})^{-1}$  exists they have the unique solution for  $\hat{C}$ ,

$$\hat{C} = (T_{m,r}^t Q_n T_{m,r})^{-1} T_{m,r}^t Q_n B. \quad (5)$$

The approximate curve by using (5) is best approximation with respect to  $L_2$  norm. The  $L_2$  best approximation is known as Legendre polynomials. See Eck[7] for detailed discussion.

For error analysis, we need the definition of the *Moore - Penrose inverse*.

DEFINITION 1. Let  $A$  be the  $n \times m$  ( $n > m$ ) matrix. The  $m \times n$  matrix  $X$  is called the *Moore - Penrose inverse* of the  $A$ , if it satisfies the following conditions known as Moore - Penrose conditions:

$$\begin{aligned} AXA &= A, & (AX)^t &= AX \\ XAX &= X, & (XA)^t &= XA \end{aligned}$$

The Moore - Penrose inverse of  $A$  is usually denoted by  $A^+$ . The Moore - Penrose inverse  $A^+$  of the matrix  $A$  is uniquely determined[13]. If  $\text{rank}(A) = m$ , then  $A^+ = (A^t A)^{-1} A^t$ , while if  $\text{rank}(A) = m = n$ , then  $A^+ = A^{-1}$ [13].

The matrix  $M = (T_{m,r}^t Q_n T_{m,r})^{-1} T_{m,r}^t Q_n$  is the Moore - Penrose inverse of the matrix  $T_{m,r}$ .

To obtain the approximation error  $\epsilon_{L_2}$ , put  $\hat{C}$  into the equation in the Problem 2.

THEOREM 3. The error of the solution  $\hat{C}$  of the Problem 2 is

$$\epsilon_{L_2}^2 = B^t Q_n [I - T_{m,r} (T_{m,r}^t Q_n T_{m,r})^{-1} T_{m,r}^t] B$$

For simple computation we may use the discrete  $l_2$  distance function, then the Problem 1 rewritten as the following:

PROBLEM 3 ( $l_2$  Degree Reduction). Find control points  $\{c_i\}_{i=0}^m$  so that the distance

$$d_{DLS}(b^n, c^m) = d_{DLS}(b^n, c^{(r)}) = \sqrt{D^t D}$$

between  $\{b_i\}_{i=0}^n$  and  $\{c_i^{(r)}\}_{i=0}^n$  is minimized.

As in the case of  $L_2$  degree reduction, we obtain the solution  $\hat{C}_{DLS}$  as

$$\hat{C}_{DLS} = (T_{m,r}^t T_{m,r})^{-1} T_{m,r}^t B$$

The matrix  $(T_{m,r}^t T_{m,r})^{-1} T_{m,r}^t$  is also the Moore – Penrose inverse of  $T_{m,r}$ . By the uniqueness of the Moore – Penrose inverse, the  $L_2$  solution  $\hat{C}$  and the discrete  $l_2$  solution  $\hat{C}_{DLS}$  are equal, i.e.

$$\hat{C} = \hat{C}_{DLS}.$$

Thus the  $L_2$  degree reduction curve and the discrete  $l_2$  degree reduction curve are same.

To obtain the approximation error  $\epsilon_{l_2}$ , put  $\hat{C}_{DLS}$  into the equation in the Problem 3.

THEOREM 4. The error of the solution  $\hat{C}_{DLS}$  of the Problem 3 is

$$\epsilon_{l_2}^2 = B^t [I - T_{m,r} (T_{m,r}^t T_{m,r})^{-1} T_{m,r}^t] B$$

The  $(m+2) * (m+2)$  matrix  $P_m = I - T_m (T_m^t T_m)^{-1} T_m^t = \{p_{i,j}\}$  has elements

$$p_{i,j} = (-1)^{i+j} \frac{\binom{m+1}{i} \binom{m+1}{j}}{\binom{2m+2}{m+1}}, i, j = 0, 1, \dots, m+1.$$

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